#### CS 201P Computer Security Winter 2020

# Process, OS Interfaces, and getting hands-on with UNIX and C

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## **OS Interfaces**

A process usually alternates between running in the user space and the kernel space Since there may be multiple user processes, the kernel uses the CPU's hardware protection mechanisms to ensure that each process executing in the user space can access only its own memory.

### fork()

```
int pid = fork();
if (pid > 0){
    printf("parent: child=%d\n", pid);
    pid = wait();
    printf('child %d is donen, pid);
   else if(pid == 0){
}
        printf("child: exiting\n);
       exit();
    else {
        printf("fork error\n");
}
```

```
int pid = fork(); // a new child process is created with same memory contents as the parent process.
if (pid > 0){
```

```
printf("parent: child=%d\n", pid);
pid = wait();
printf('child %d is done\n", pid);
else if(pid == 0){
    printf("child: exiting\n);
    exit();
else {
    printf("fork errorn");
```

int pid = fork(); \* // a new child process is created with same memory contents as the parent process. if (pid > 0){

```
printf("parent: child=%d\n", pid);
pid = wait ();
printf('child %d is done\n", pid);
else if(pid == 0){
    printf("child: exiting\n);
    exit();
else {
```

```
printf("fork error\n");
```

At this point we have two processes running.

> fork returns 0 to the child process.

> fork returns the pid of the child process in the parent process.

```
int pid = fork(); // a new child process is created with same memory contents as the parent process.
if (pid > 0){ // the parent process
```

```
printf("parent: child=%d\n", pid);
pid = wait ();
printf('child %d is done\n", pid);
```

else {

```
printf("fork error\n");
```

```
int pid = fork(); // a new child process is created with same memory contents as the parent process.
if (pid > 0){ // the parent process
    printf("parent: child=%d\n", pid);
    pid = wait (); // returns the id of an exited child process
```

printf('child %d is done\n", pid);

} else if(pid == 0){ // the child process
 printf("child: exiting\n);
 exit();

else {

printf("fork error\n");

```
int pid = fork(); \quad // a new child process is created with same memory contents as the parent process.if (pid > 0) \{ // the parent process
```

```
printf("parent: child=%dn", pid);
pid = wait (); // returns the id of an exited child process
printf('child %d is donen, pid);
else if(pid == 0){ // the child process
     printf("child: exiting\n);
     exit(); // exits the calling process, releasing resources
else {
     printf("fork errorn");
```

So what's the output? Say child pid = 1234.

```
int pid = fork(); \quad // a new child process is created with same memory contents as the parent process.if (pid > 0) \{ // the parent process
```

```
printf("parent: child=%dn", pid);
pid = wait (); // returns the id of an exited child process
printf('child %d is donen, pid);
else if(pid == 0){ // the child process
     printf("child: exiting\n);
     exit(); // exits the calling process, releasing resources
else {
     printf("fork errorn");
```

int pid = fork(); // a new child process is created with same memory contents as the parent process. if (pid > 0){ // the parent process printf("parent: child=%d\n", pid); pid = wait (); // returns the id of an exited child process printf('child %d is donen, pid); Output: else if(pid == 0){ // the child process parent: child=1234 child: exiting printf("child: exiting\n); parent: child 1234 is done exit(); // exits the calling process, releasing resources OR else { child: exiting printf("fork errorn"); parent: child=1234 parent: child 1234 is done

While the memory contents of the child and parent processes are *initially* the same, changing a variable in one does not affect the other.

#### exec()

exec() replaces the calling process's memory with a new memory image loaded from a file stored in the file system. When exec succeeds, it does not return to the calling program.

Ref:

xv6 - a simple, Unix-like teaching operating system, Cox et al., PDOS, CSAIL, MIT