

CS 201P Computer Security Winter 2020

The buffer overflow attack: Exploring the stack and the attack strategy

7 February 2020
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the playfield

S

H

BSS

DS

TS

buffer

S

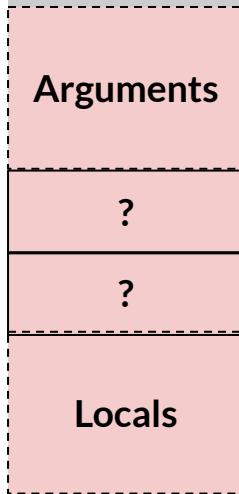
H

BSS

DS

TS

buffer



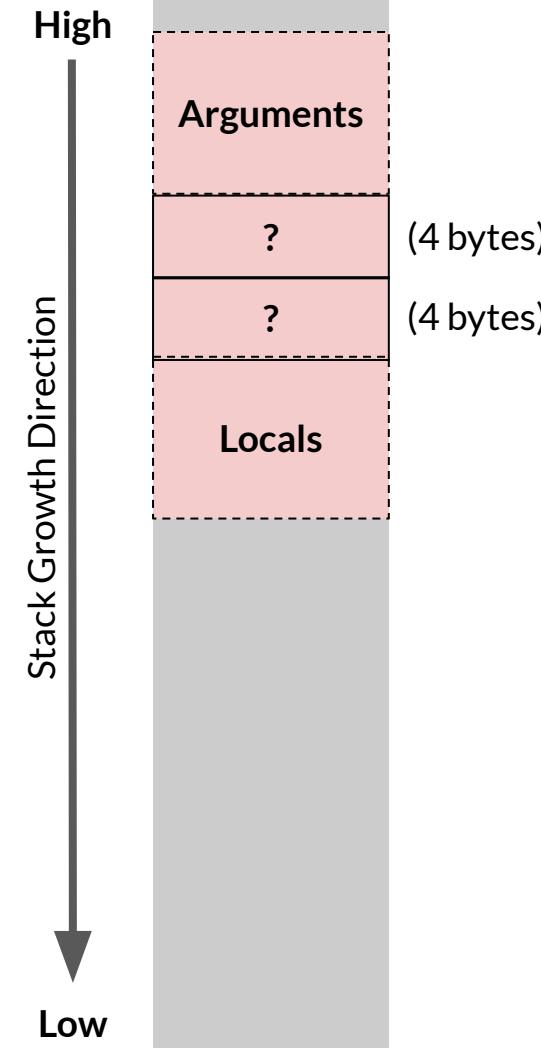
stack

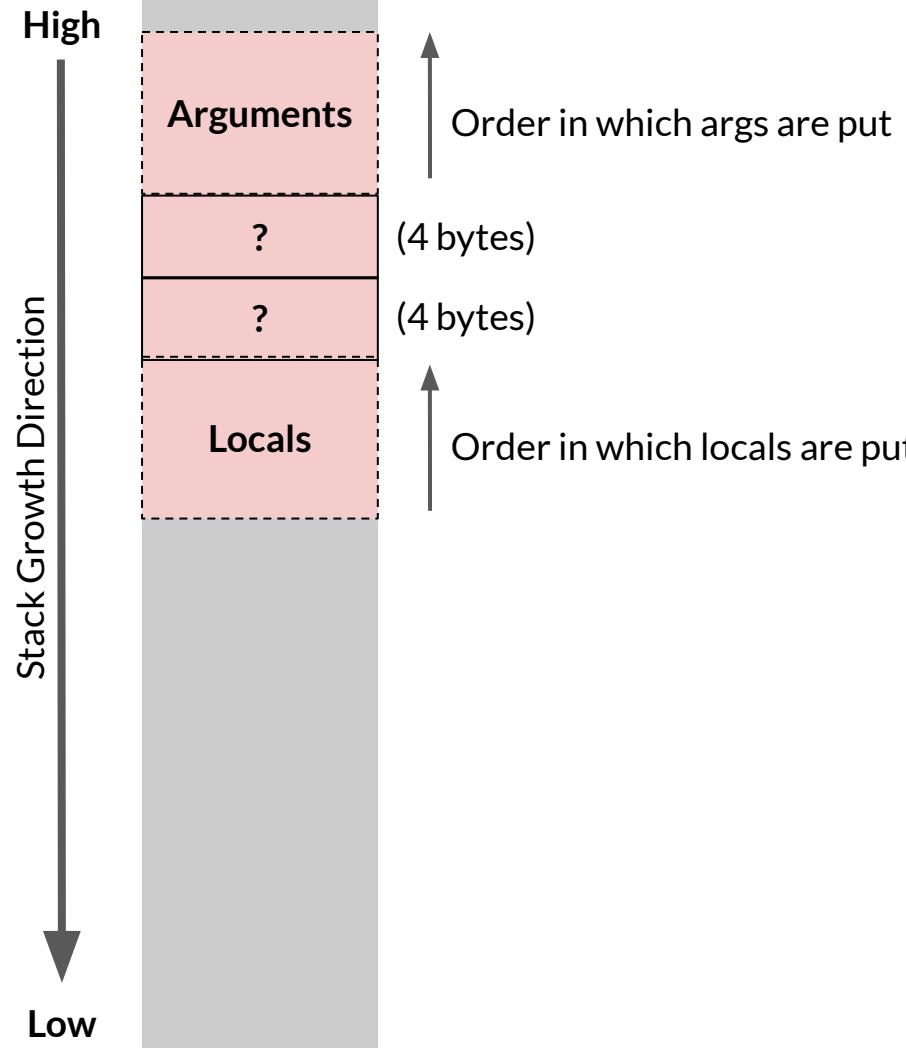
(32-bit architecture)

after any function invocation

stack

(32-bit architecture)





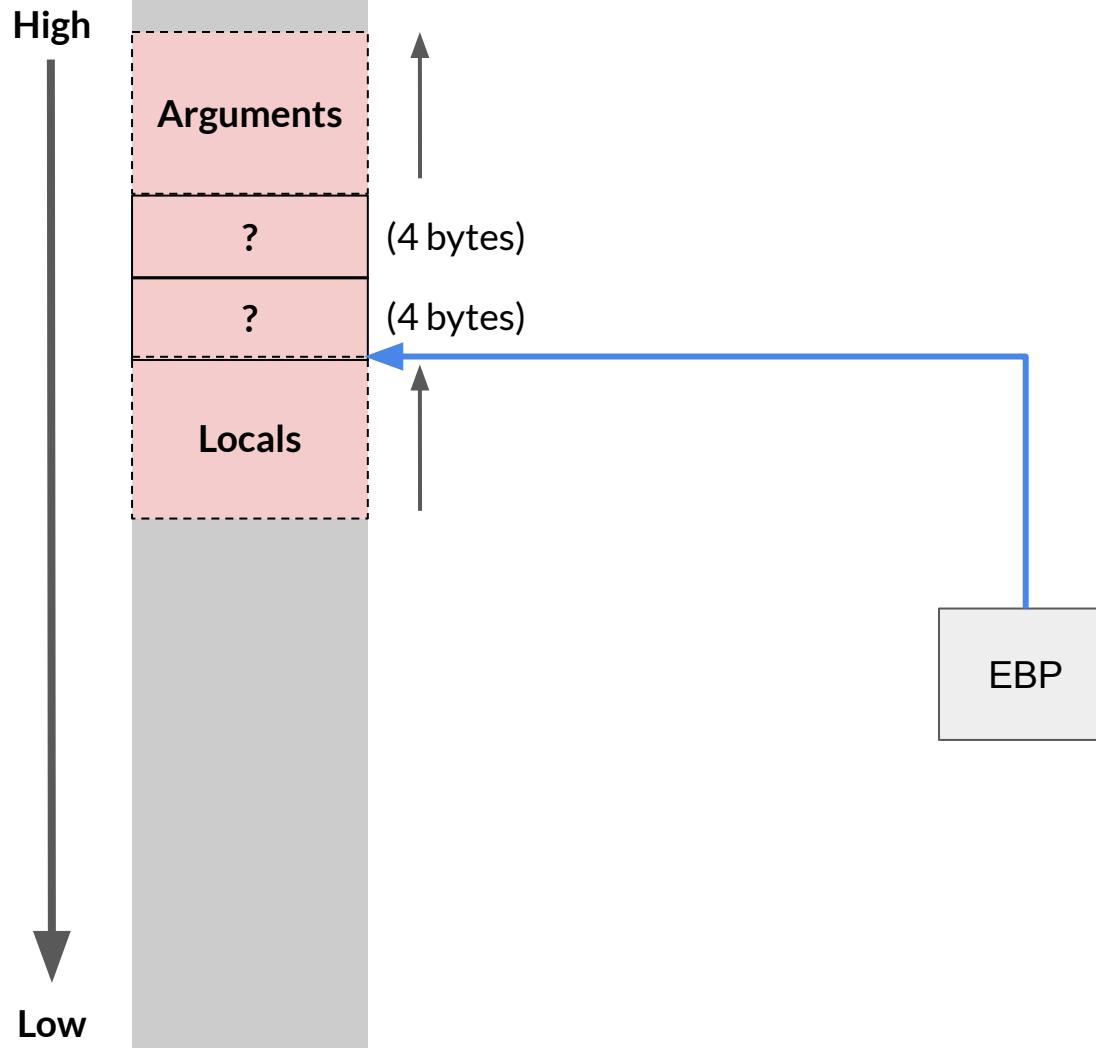
stack

(32-bit architecture)

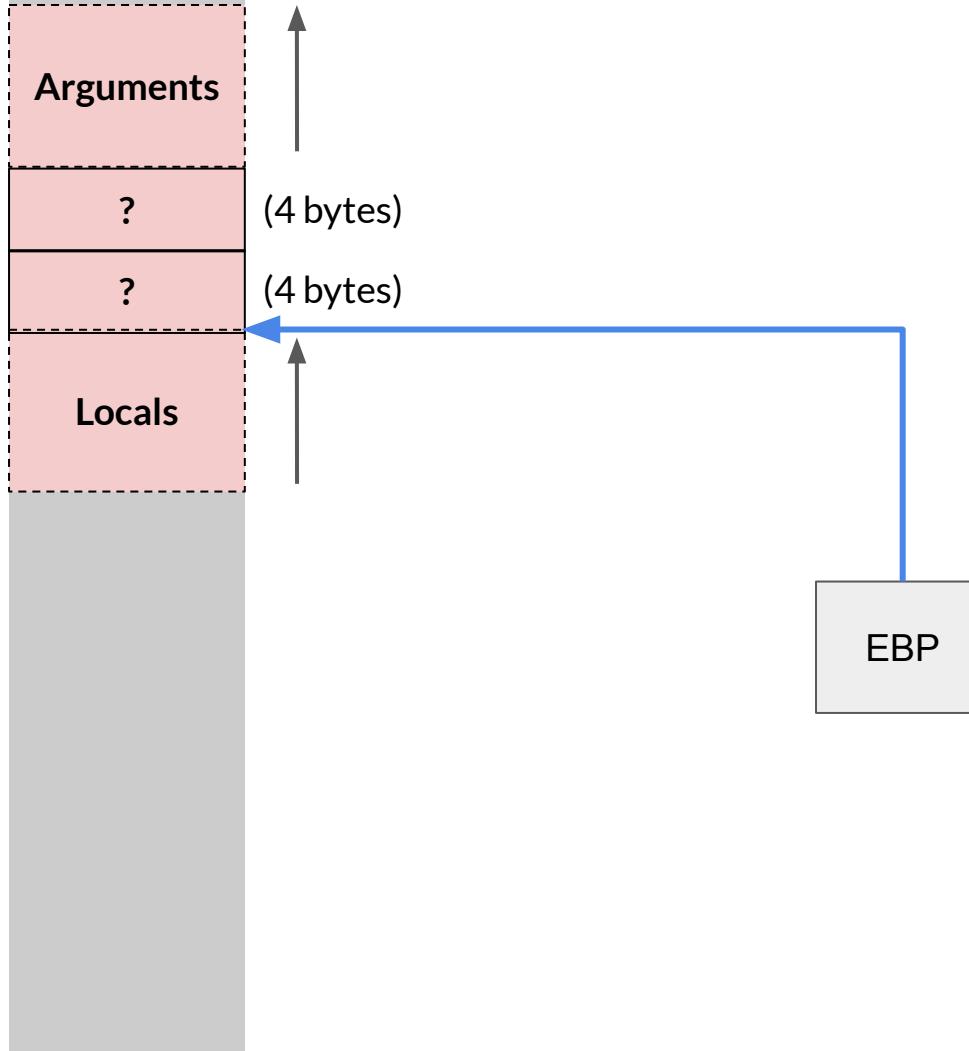
*how to refer to
the variables?*

stack

(32-bit architecture)



High



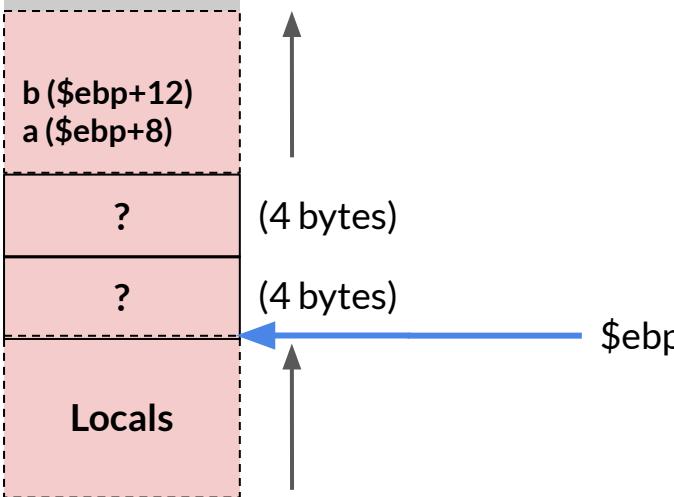
stack

(32-bit architecture)

```
int main(int a, int b){
```

...

High



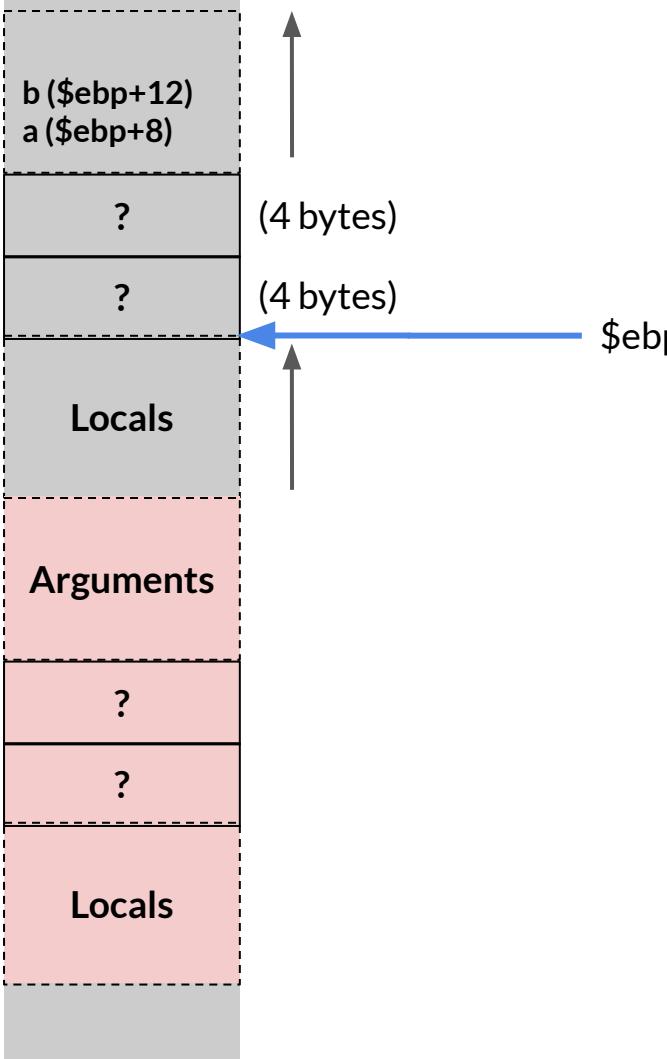
stack

(32-bit architecture)

our function calls another function

Low

High



stack

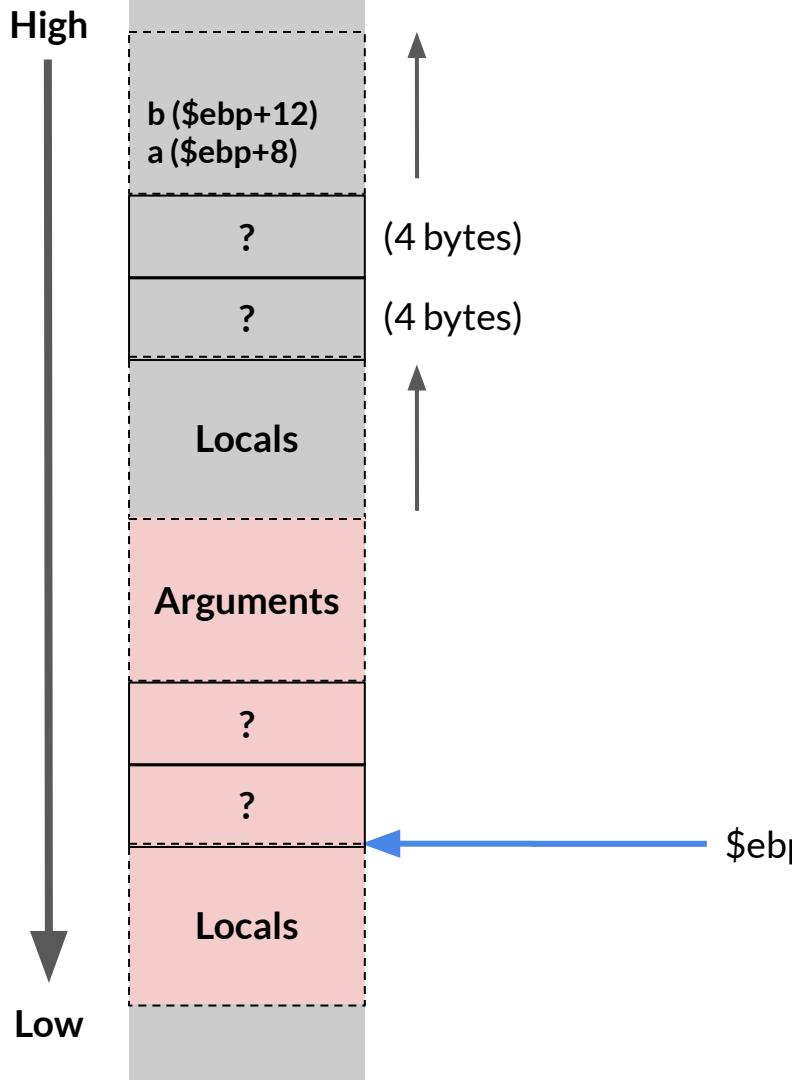
(32-bit architecture)

our function calls another function

stack

(32-bit architecture)

our function calls another function

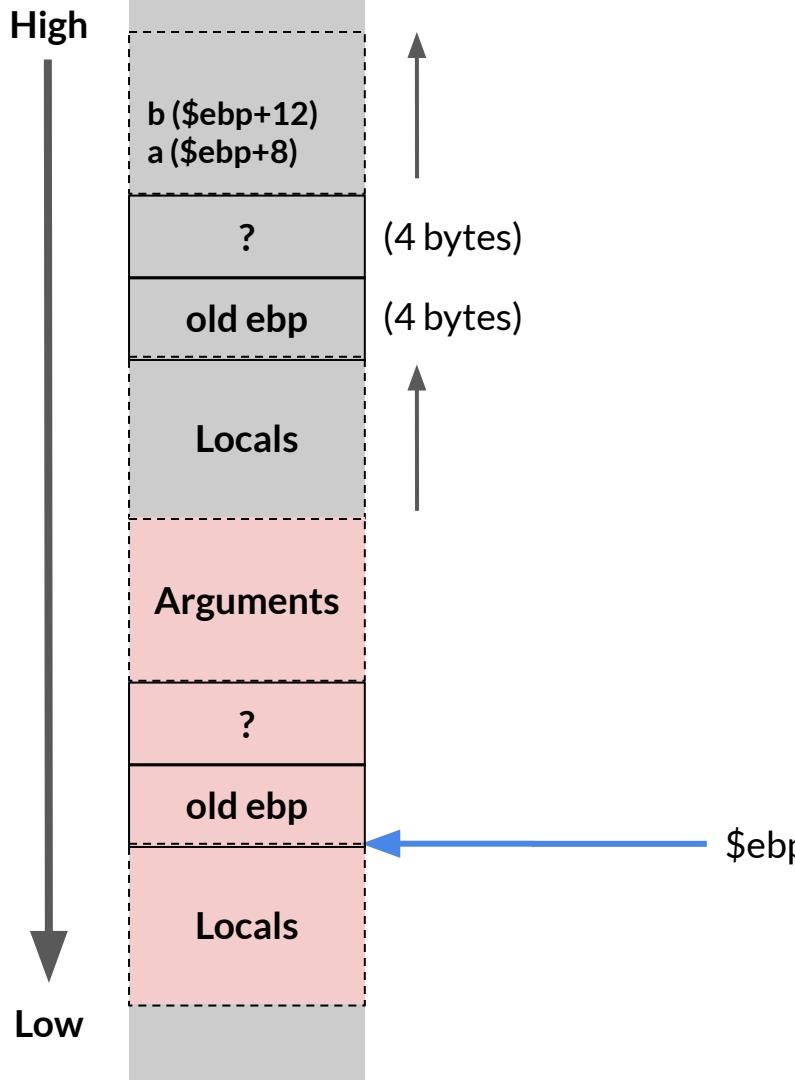


*The old \$ebp
value?*

stack

(32-bit architecture)

our function calls another function



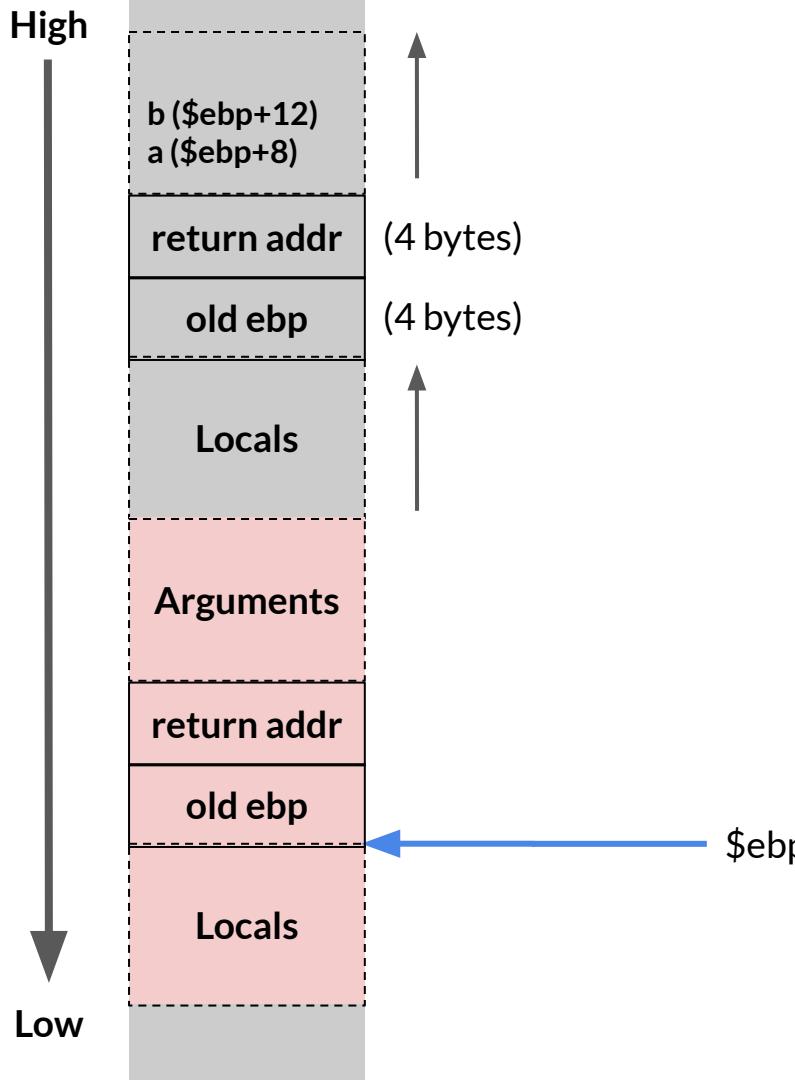
*Which instruction do we
execute after we are done
with the current function?*

*The instruction right after
the instruction that called
the function.*

stack

(32-bit architecture)

our function calls another function

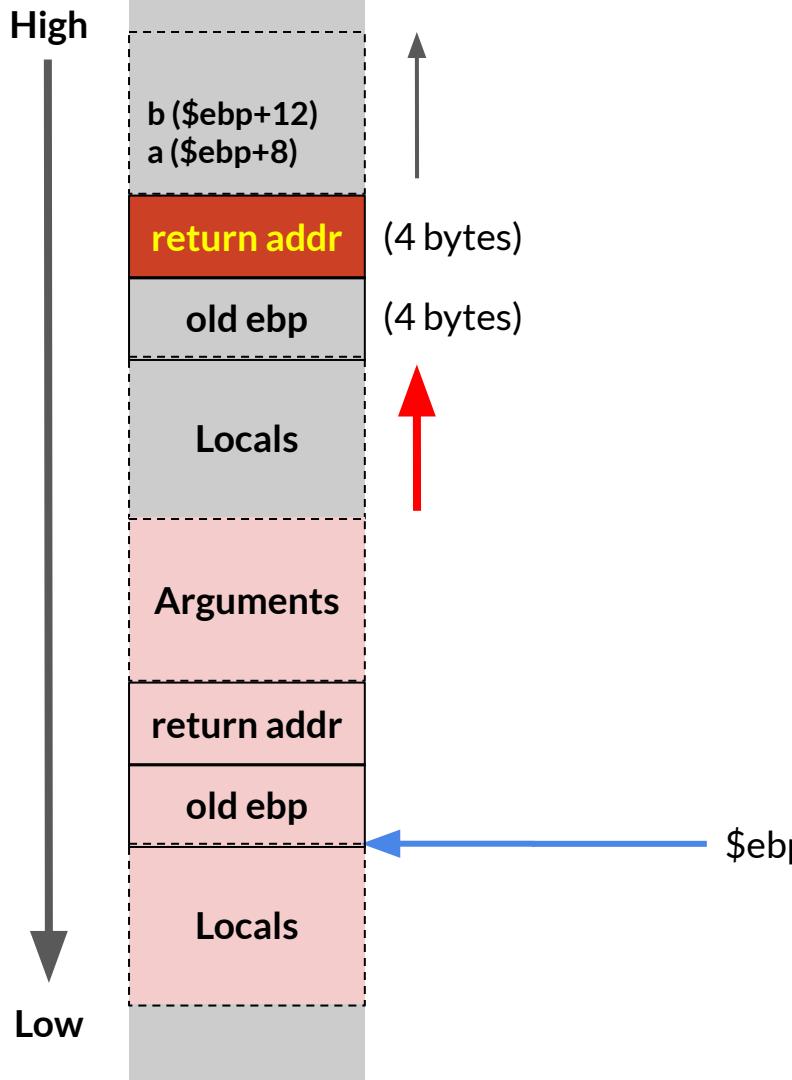


So what's interesting?

stack

(32-bit architecture)

our function calls another function



*If we can make return
address point to our code...*

*But we need to be careful of
what we put in return
addresss...*

our strategy

taking aim
(address calculation)

Vulnerable Program

Arguments

return addr

old ebp

Locals

```
int bof(char *str)
{
    char buffer[24];
    strcpy(buffer, str);
    return 1;
}

int main(int argc, char **argv)
{
    char str[517];
    FILE *badfile;
    badfile = fopen("badfile", "r");
    fread(str, sizeof(char), 517, badfile);
    bof(str);
    printf("Returned Properly\n");
    return 1;
}
```

Vulnerable Program

Arguments

return addr

old ebp

Locals

```
int bof(char *str)
{
    char buffer[24];
strcpy(buffer, str);
    return 1;
}

int main(int argc, char **argv)
{
    char str[517];
    FILE *badfile;
    badfile = fopen("badfile", "r");
    fread(str, sizeof(char), 517, badfile);
    bof(str);
    printf("Returned Properly\n");
    return 1;
}
```

Vulnerable Program

Arguments

return addr

old ebp

Locals

```
int bof(char *str)
{
    char buffer[24];
strcpy(buffer, str); //loads your exploit code to the memory
    return 1;
}

int main(int argc, char **argv)
{
    char str[517];
    FILE *badfile;
    badfile = fopen("badfile", "r");
    fread(str, sizeof(char), 517, badfile);
    bof(str);
    printf("Returned Properly\n");
    return 1;
}
```

Vulnerable Program

```
int bof(char *str)
{
    char buffer[24];
    strcpy(buffer, str); //loads your exploit code to the memory
    return 1;              //just need to set up return address of this fn
}                          //after we are done with the strcpy instruction
```

```
int main(int argc, char **argv)
{
    char str[517];
    FILE *badfile;
    badfile = fopen("badfile", "r");
    fread(str, sizeof(char), 517, badfile);
    bof(str);
    printf("Returned Properly\n");
    return 1;
}
```

Arguments

return addr

old ebp

Locals

```
$gcc -z execstack -fno-stack-protector -g -o stack_dbg stack.c  
$touch badfile
```

```
$gdb -q stack_dbg
```

Enter GDB

bof() stack

Arguments

return addr

old ebp

Locals

```
$gcc -z execstack -fno-stack-protector -g -o stack_dbg stack.c  
$touch badfile
```

```
$gdb -q stack_dbg
```

Enter GDB

```
break at bof() function  
gdb$ b bof()  
gdb$ run
```

bof() stack

Arguments

return addr

old ebp

Locals

```
$gcc -z execstack -fno-stack-protector -g -o stack_dbg stack.c  
$touch badfile
```

```
$gdb -q stack_dbg
```

Enter GDB

```
break at bof() function
```

```
gdb$ b bof()
```

```
gdb$ run
```

bof() stack

Arguments

return addr

old ebp

\$buffer

The program will stop after bof() is entered.

Find the location of the buffer variable.

```
gdb$ p/x &buffer
```

[you get the hex address of buffer]

```
$gcc -z execstack -fno-stack-protector -g -o stack_dbg stack.c  
$touch badfile
```

```
$gdb -q stack_dbg
```

Enter GDB

```
break at bof() function
```

```
gdb$ b bof()
```

```
gdb$ run
```

bof() stack

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return addr

old ebp

\$buffer

The program will stop after bof() is entered.

Find the location of the buffer variable.

```
gdb$ p/x &buffer
```

[you get the hex address of buffer]

Get address of \$ebp value

```
gdb$ p/d addr(ebp) - addr(buffer)
```

```
$gcc -z execstack -fno-stack-protector -g -o stack_dbg stack.c  
$touch badfile
```

```
$gdb -q stack_dbg
```

Enter GDB

```
break at bof() function
```

```
gdb$ b bof()
```

```
gdb$ run
```

bof() stack

Arguments

return addr

old ebp

\$buffer

The program will stop after bof() is entered.
Find the location of the buffer variable.

```
gdb$ p/x &buffer
```

[you get the hex address of buffer]

Get address of \$ebp value

108

```
gdb$ p/d addr(ebp) - addr(buffer)
```

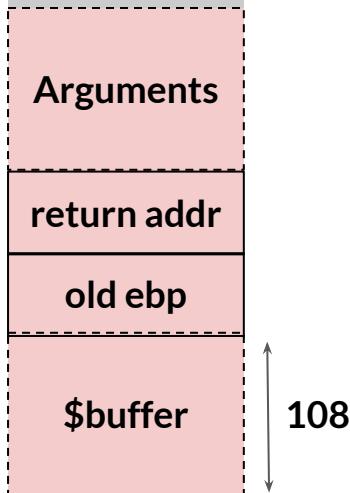
Let's say it's 108

Use the offsets you get to set up the return address value.

NOTE:

Regarding the content of the return address, as discussed, it would be \$ebp + [some value, like 12]. Make sure the resulting hexadecimal address from the sum does not contain any sequence like “00”: This sequence would resemble a null byte. When using strcpy, the copy process would stop on encountering the null byte, causing the attack to fail.

bof() stack



Thank you

Ref. Computer Security: A Hands on Approach
Wenliang Du