

238P Operating Systems, Spring 2019

Topics : Paging and Stacks

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Aftab Hussain

University of California, Irvine

Review of Address Translation using Paging

```
mov (%EBX), EAX # mov value from the location pointed by EBX into EAX  
EAX = 0  
EBX = 20 983 809
```

20 983 809 =

00 0000 0101	00 0000 0011	0000 0000 0001
--------------	--------------	----------------

- > With paging enabled, xv6 can deal with 32 bit virtual addresses
- > Size of the virtual address space is 2^{32} bytes = 4GB
- > A page is of size 4KB
- > There are ~1 million pages in the virtual address space.
- > The pages are mapped to a physical address space, limited by the physical memory.
- > We shall take the help of two tables (2-level tree) to map from virtual to physical.
- > These tables are stored in the physical memory.

mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

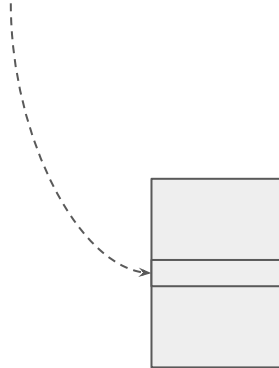
EAX = 0

EBX = 20 983 809

20 983 809 =

00 0000 0101	00 0000 0011	0000 0000 0001
--------------	--------------	----------------

points to
entry
in the 1st
table
(the PD)



page directory
(1024 entries)

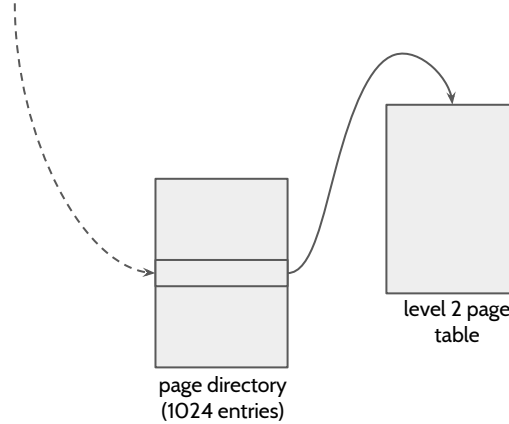
mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

EAX = 0

EBX = 20 983 809

20 983 809 = 00 0000 0101 | 00 0000 0011 | 0000 0000 0001

points to
entry
in the 1st
table
(the PD)



mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

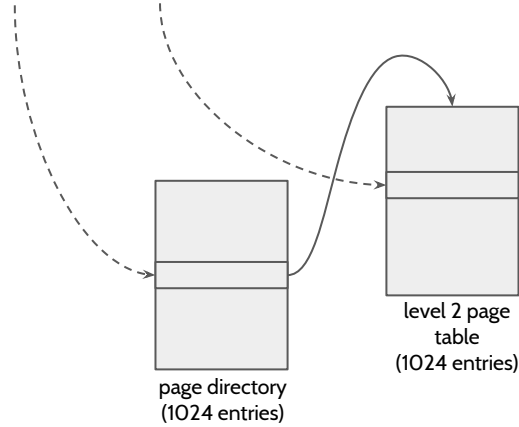
EAX = 0

EBX = 20 983 809

20 983 809 = 00 0000 0101 | 00 0000 0011 | 0000 0000 0001

points to
entry
in the 1st
table
(the PD)

points to
entry
in the 2nd
table



mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

EAX = 0

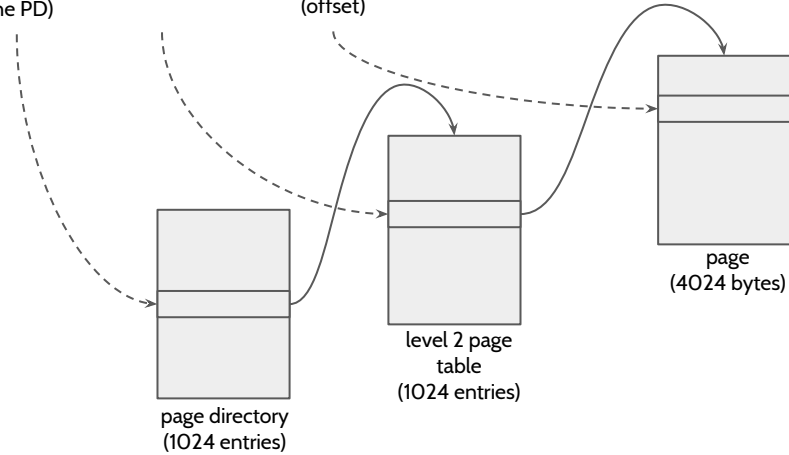
EBX = 20 983 809

20 983 809 = 00 0000 0101 00 0000 0011 0000 0000 0001

points to
entry
in the 1st
table
(the PD)

points to
entry
in the 2nd
table


points to
the actual
byte in the
page
(offset)




```
mov (%EBX), EAX # mov value from the location pointed by EBX into EAX  
EAX = 0  
EBX = 20 983 809
```

20 983 809 =

00 0000 0101	00 0000 0011	0000 0000 0001
--------------	--------------	----------------


page number

mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

EAX = 0

EBX = 20 983 809

20 983 809 = 00 0000 0101 00 0000 0011 0000 0000 0001

page number

1M (1,048,575)

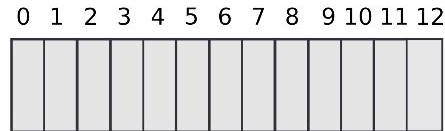
Virtual Address
Space (or Memory)
of the Process



0 1 2

page number = 5123
or (0b1 0100 0000 0011)

Physical
Memory



mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

EAX = 0

EBX = 20 983 809

20 983 809 = 00 0000 0101 00 0000 0011 0000 0000 0001

page number

1M (1,048,575)

Virtual Address
Space (or Memory)
of the Process



CR3 = 0

0 1 2

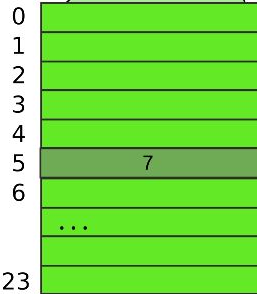
page number = 5123
or (0b1 0100 0000 0011)

0 1 2 3 4 5 6 7 8 9 10 11 12

Physical
Memory



32 bits (4 bytes)



Level 1
(Page Table
Directory)

mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

EAX = 0

EBX = 20 983 809

20 983 809 = 00 0000 010 00 0000 0011 0000 0000 0001

page number

1M (1,048,575)

Virtual Address
Space (or Memory)
of the Process



CR3 = 0

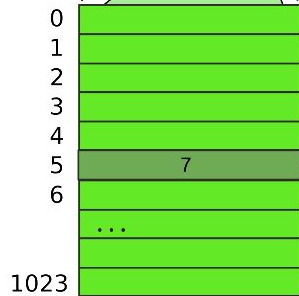
0 1 2

page number = 5123
or (0b1 0100 0000 0011)

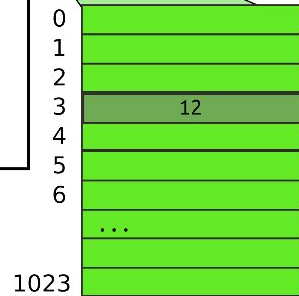
0 1 2 3 4 5 6 7 8 9 10 11 12

Physical
Memory

32 bits (4 bytes)



Level 1
(Page Table
Directory)



Level 2
(Page Table)

mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

EAX = 0

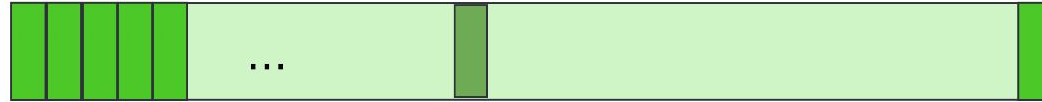
EBX = 20 983 809

20 983 809 = 00 0000 0101 00 0000 0011 0000 0000 0001

page number

1M (1,048,575)

Virtual Address
Space (or Memory)
of the Process



CR3 = 0

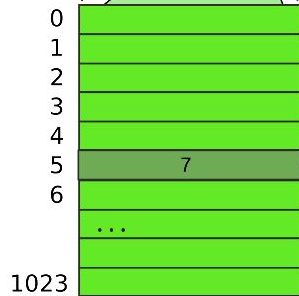
0 1 2

page number = 5123
or (0b1 0100 0000 0011)

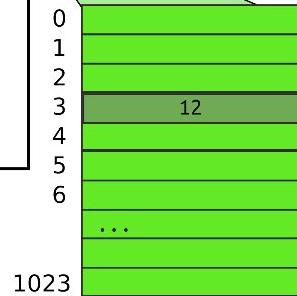
0 1 2 3 4 5 6 7 8 9 10 11 12

Physical
Memory

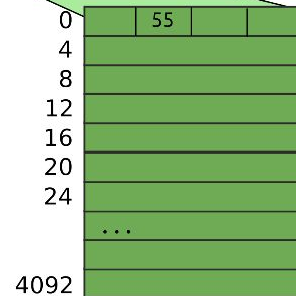
32 bits (4 bytes)



Level 1
(Page Table
Directory)



Level 2
(Page Table)



Page

mov (%EBX), EAX # mov value from the location pointed by EBX into EAX

EAX = 0

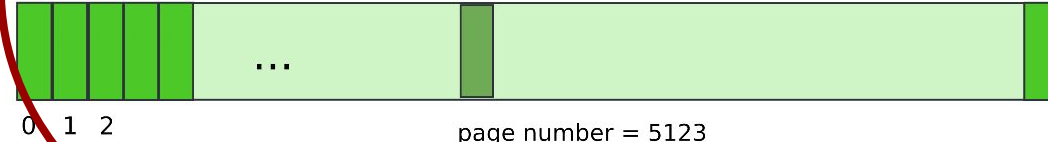
EBX = 20 983 809

20 983 809 = 00 0000 0101 00 0000 001 0000 0000 0001

page number

1M (1,048,575)

Virtual Address
Space (or Memory)
of the Process

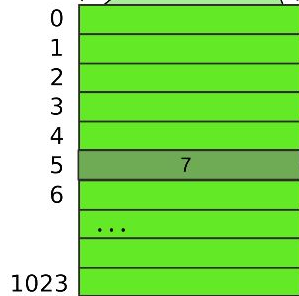


CR3 = 0

0 1 2 3 4 5 6 7 8 9 10 11 12

Physical
Memory

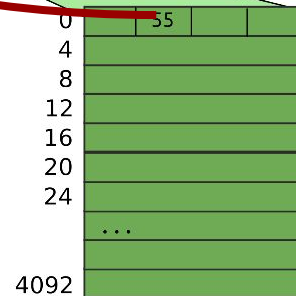
32 bits (4 bytes)



Level 1
(Page Table
Directory)



Level 2
(Page Table)



Page

- > Problem 2, from [CS238P Fall 2018, Midterm](#)
- > Problem 1, from [CS238P Winter 2018, Midterm](#)

Stack

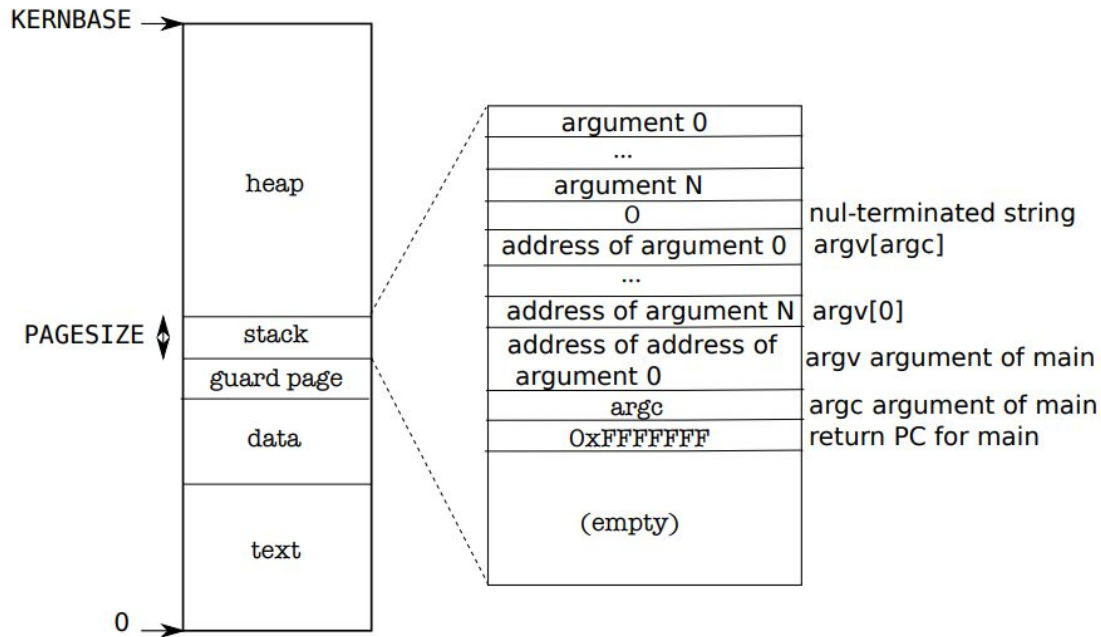


Figure 2-3. Memory layout of a user process with its initial stack.

- > The stack is populated from the higher to the lower addresses.
- > Arguments of function calls occupy the higher addresses of the stack.
- > The guard page is unmapped, and is aimed to prevent the stack from growing beyond the size of the stack (which is 1 page.)

- > Problems 2 and 3 from [CS238P Winter 2018, Midterm](#)
- > Problem 3, from [CS238P Fall 2018, Midterm](#)